

*ECRTS 2012 in Pisa, Italy*  
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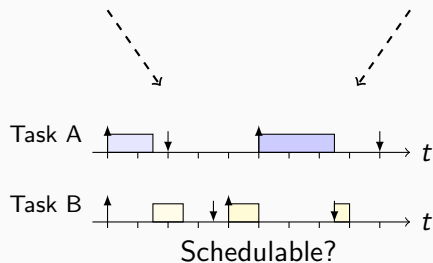
## Hardness Results for Static Priority Real-Time Scheduling

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# Problem Overview



**This Work:**  
Test Complexity

# Model/Design Choices

## Task Models:

- Periodic (L&L)
- Generalized Multiframe (GMF)
- Digraph Real-Time (DRT)
- ...

## Schedulers:

- Dynamic Priorities: EDF
- Static Priorities

Different combinations = Different complexity

(Here: Uniprocessor, preemption, precise tests)

# Complexity of Schedulability Test

- Efficient schedulability tests possible?
  - ▶ (“Efficient” = “pseudo-polynomial”)

	EDF	Static
L&L	Yes	Yes
GMF	Yes	Yes* No!
DRT	Yes	No
EDRT	No	No

- \* = Takada & Sakamura, 1997
- *Flawed!*

## Theorem (Our technical result)

For GMF task systems, the schedulability problem for *static* priority schedulers is *strongly coNP-hard*.

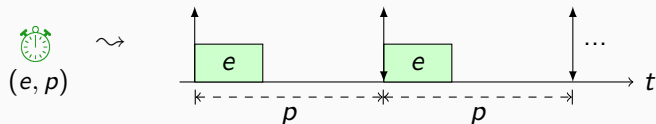
# Fahrplan

- 1 Problem Overview
- 2 Task Models: L&L, GMF, DRT
- 3 Analysis Methods
  - EDF: Demand Bound Function
  - Static Priorities: Maximum Interference Function
- 4 Hardness Result

# The Liu and Layland (L&L) Task Model

(Liu and Layland, 1973)

- Tasks are *periodic*
  - ▶ Job WCET  $e$
  - ▶ Minimum inter-release delay  $p$  (implicit deadline)

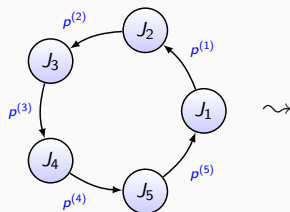


- Advantages: Well-known model; efficient schedulability tests
- However, *not everything is periodic...*

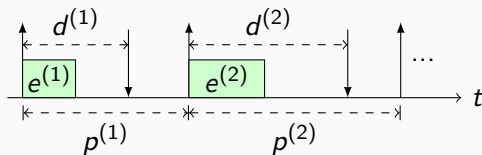
# The General Multiframe (GMF) Task Model

(Baruah et al, 1999)

- Tasks *cycle* through job types, “frames”
  - ▶ Vector for WCET ( $e^{(1)}, \dots, e^{(n)}$ )
  - ▶ Vector for deadlines ( $d^{(1)}, \dots, e^{(n)}$ )
  - ▶ Vector for minimum inter-release delays ( $p^{(1)}, \dots, p^{(n)}$ )



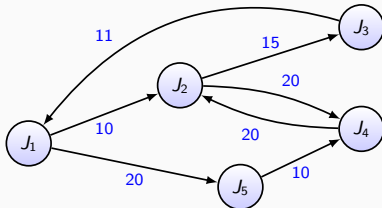
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# The Digraph Real-Time (DRT) Task Model

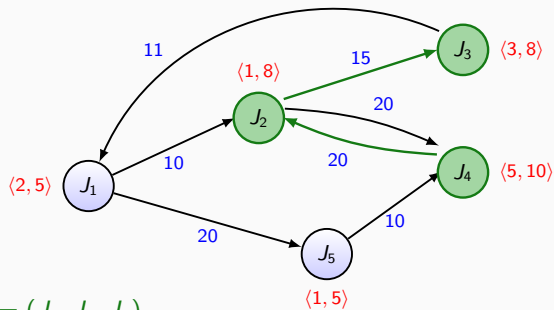
(S. et al, RTAS 2011)

- Branching, cycles (loops), ...
- *Directed graph* for each task
  - ▶ Vertices  $J$ : jobs to be released (with WCET and deadline)
  - ▶ Edges  $(J_i, J_j)$ : minimum inter-release delays  $p(J_i, J_j)$

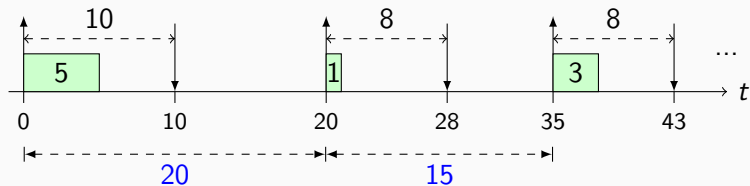




# DRT: Semantics



Path  $\pi = (J_4, J_2, J_3)$



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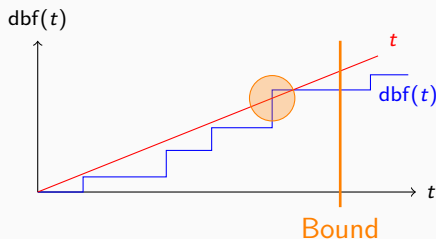
# The Demand Bound Function

- Given a time interval length  $t$
- $\text{dbf}(t)$  bounds the *demand* for processor time within *any*  $t$  interval

## Theorem

A task system  $\tau$  is schedulable with EDF iff:

$$\forall t \geq 0 : \sum_{T \in \tau} \text{dbf}_T(t) \leq t$$



# Complexity of Schedulability Test

- Efficient schedulability tests:

	EDF	Static
L&L	Yes	Yes
GMF	Yes	?
DRT	Yes	
EDRT	No	

# Schedulability for Static Priorities

- L&L tasks: *Response Time Analysis*

$$R_i = C_i + \underbrace{\sum_{j \in hp(i)} \left\lceil \frac{R_i}{T_j} \right\rceil \cdot C_j}_{\text{Interference Term}}$$

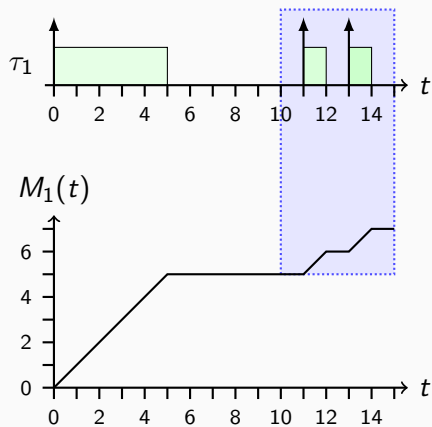
- Generalize for GMF: *Maximum Interference Function (MIF)*
  - ▶  $M_j(t)$ : Maximum interference that  $\tau_j$  can cause within  $t$  time units

$$R_i = C_i + \sum_{j \in hp(i)} M_j(R_i)$$

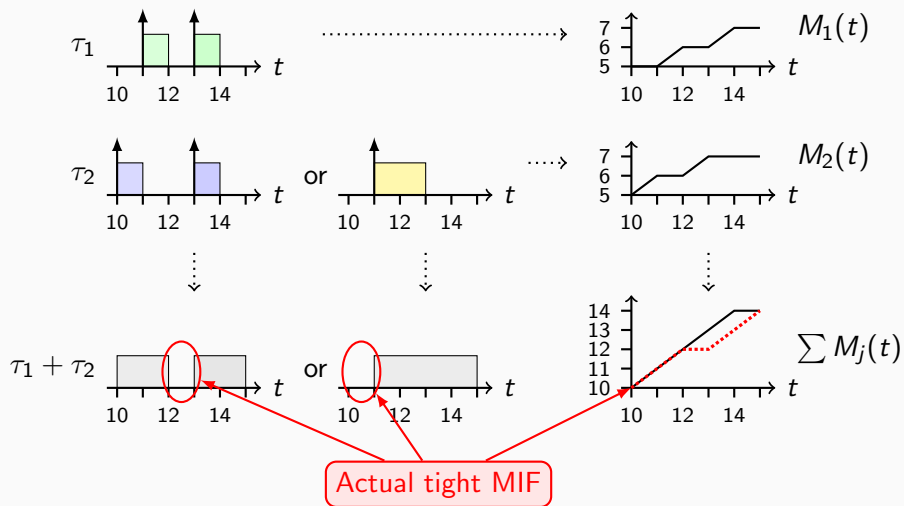
- ▶ Efficiently computable for GMF/DRT/...

**BUT: Inherently overapproximate!**

# MIF: Example



# MIF: Combined Example



# Schedulability for Static Priorities

- In summary: MIF is *pessimistic*
- Possible improvement?
  - ▶ Define MIF additions better? Precise?
  - ▶ Use another abstraction level?
  - ▶ ...?
- No!

## Theorem (Our technical result)

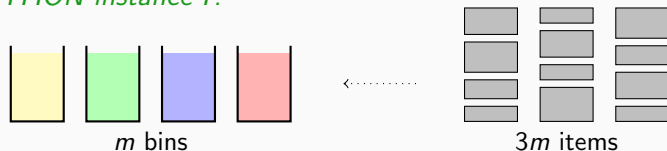
*For GMF task systems, the schedulability problem for **static** priority schedulers is **strongly coNP-hard**.*

- Thus: No precise efficient analysis possible.



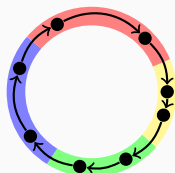
# Hardness Result: Proof sketch

*3-PARTITION instance I:*



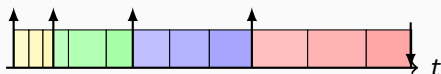
Possible to (exactly) fit all items? (strongly NP-hard)

*Reduction to GMF schedulability:*



One GMF task  
for *each* item

Packing possible  $\iff$  Busy interval



Thus:  $\tau(I)$  unshed.  $\iff I \in 3\text{-PARTITION}$

# Summary and Outlook

	EDF	Static
L&L	Yes	Yes
GMF	Yes	Yes No!
DRT	Yes	No
EDRT	No	No

- Showed *intractability* of static scheduling for GMF
- Insight:
  - ▶ For EDF, “simple” overload test suffices (Local worst cases combine.)
  - ▶ For static prio: More structure → more complex test (Local worst case unclear.)
- Ongoing work:
  - ▶ Solve anyway? Heuristics?
  - ▶ Use SAT-/SMT-solvers

Thanks!